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Two very efficient algorithms for generating pseudorandom numbers from the gamma distribution have been developed by Ahrens and Dieter; in the present work these are combined with a third method to produce a combination generator capable of excellent performance for any order of gamma variate. The algorithms are briefly described and an IBM 360 Assembler implementation of them is described and tested. A second computer program for the generation. of pseudorandom Cauchy deviates is presented; this program uses a new

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20. (continued) algorithm which is also described. Both computer programs are intended to be used with the Naval Postgraduata School random number package LLRANDOM.

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# GENERATING GAMMA AND CAUCHY RANDOM VARIABLES: AN EXTENSION TO THE NAVAL POSTGRADUATE SCHOOL RANDOM NUMBER PACKAGE

by

D. W. Robinson and P. A. W. Lewis \*

\* Work partially supported by the National Science Foundation under grant AG 476.

# NONUNIFORM RANDOM NUMBER PACKAGE

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# I. Introduction

The use of uniformly or non-uniformly distributed pseudorandom numbers in systems simulation, statistical sampling experiments and analytical Monte Carlo work is by now well established. Numerous algorithms exist for producing such numbers from various distributions; for summaries of common techniques, see Knuth [5], Gaver and Thompson [2] or Ahrens and Dieter [1].

The user of pseudorandom numbers is usually concerned with the details of the algorithm employed but rather with the results; a good algorithm, then, which is fast, uses minimum computer memory and produces numbers with satisfactory statistical properties. The for statistically competent algorithms for search pseudorandom numbers has resulted in the specification of many so-called "exact" generators, that is those whose deviation from the true distribution concerned is the result of computer rounding errors rather than any defect in the method itself. Such methods for nonuniform random numbers are often based on the assumption that "good" uniform numbers are available from an independent generator.

Exact generators for nonuniform pseudorandom numbers are often quite complex and so assembly-level coding is often resorted to when implementing them in order to meet the computer time and memory constraints on a good algorithm. An example is the LLRANDOM package developed at the Naval Postgraduate School by G.P. Learmonth and P.A.W. Lewis and described in [7]; it produces pseudorandom numbers

from uniform, normal and exponential distributions. This report describes an extension to the LLRANDOM package for Cauchy and gamma distributed numbers.

The Cauchy distribution has density function

(1) 
$$f(x) = \frac{1}{\pi} - \frac{1}{1 - x^2}, -\infty < x < \infty$$

and distribution function

$$F(x) = \frac{1}{2} + \frac{1}{\pi} \tan^{-1} x$$
.

While the shape of the Cauchy density resembles the normal density, the tails are much heavier; in fact, Cauchy 121... variables have no expectation and an infinite variance. The density has mode at zero and often in applications the variates are often shifted by a location parameter T or scaled by multiplying by a scale parameter S. Because of the heavy tails, Cauchy variates might find application as a "pathological" case in a systems simulation study as well as in statistical sampling experiments for robust estimation techniques. See Chapter 16 of Johnson and Kotz [4] for further details on the Cauchy distribution.

The <u>qamma distribution</u> with shape parameter A and scale parameter s has the density function

(2) 
$$f(x) = \frac{A}{S} \frac{A-1}{\Gamma(A)} - sx$$

where F(A) is Euler's gamma function

(3) 
$$\Gamma(\lambda) = \int_0^\infty x^{\lambda-1} e^{-x} dx .$$

Note that  $\Gamma(n) = (n-1)!$  when n is a non-negative integer. If the random variable X has density (2) then

$$E[X] = A / s$$

#### $V[X] = \lambda / s^2 .$

A;

When  $\lambda = 1$ , X has the exponential distribution while X, suitably scaled, has an asymptotically normal distribution as  $\lambda \rightarrow \infty$ .

We note that if X has a  $\Gamma(\lambda,1)$  distribution then X/s has a  $\Gamma(\lambda,s)$  distribution, so we may set s=1 in (2) as far as the generating algorithm is concerned. The output from the generator may then be appropriately scaled.

Gamma random variables are used in a wide variety of applications: for analytical modeling, in reliability theory and for statistical testing (the chi-squared random variable with n degrees of freedom has the  $\Gamma(n, 1)$  distribution). See [6] or Chapter 17 of [4] for more details.

## II. Use of the Subroutines

This extension to LLRANDOM is composed of two independent IBM System/360 Assembler-coded subroutines: CAUCHY for Cauchy-distributed variates and GAMA for gamma variates. The name GAMA was chosen so as not to conflict with the IBM mathematical library subprogram GAMMA which computes the gamma function (3).

The basic conventions for using GAMA and CAUCHY are the same as in the LLRANDOM package: the invoking statements

CALL CAUCHY ( IX, X, N ) and CALL GAMA ( A, IX, X, N )

will result in a vector X(1), ..., X(N) of Cauchy or  $\Gamma$  (A,1.0) pseudorandom variates, respectively. The argument IX is, in both cases, an <u>integer</u> seed to be used in the multiplicative congruential uniform generator employed by LLRANDOM. IX should be initialized just once in the calling program to some positive integer value and should <u>not</u> be altered thereafter.

The subroutine GAMA requires a source for normal and exponential deviates; these are obtained directly from the LLRANDOM package and so the statement "CALL OVFLOW" must appear once in the calling program to initialize LLRANDOM. As mentioned previously, the output from GAMA must be scaled if the scale parameter is other than one; the following set of statements will thus be required to generate a vector of 100 chi-squared variates with seven degrees of freedom:

DIMENSION X (100)

CALL OVPLOW

IX = 13726

CALL GAMA ( 3.5, IX, X, 100 )

DO 50 I = 1,100 X(I) = 2.0 + X(I) 50 CONTINUE

END

Cauchy variates are also often modified by location and scale parameters; since no expectations exist, however, we cannot refer to these parameters in terms of mean or variance. Subroutine CAUCHY is completely independent of LLRANDOM or any other subroutines so that the "CALL OVFLOW" statement is not necessary in this case. To use CAUCHY to produce a single variate C with location parameter T and scale parameter S we may use the statements

IX = 217663541
...

CALL CAUCHY ( IX, C, 1 )

C = S \* C + T
...

END

Just as in LLRANDOM, linkage overhead between the calling program and GAMA or CAUCHY will be minimized if a vector of several variates is obtained at the same time instead of just a single one. The gain in this case can be as much as 50 microseconds per variate in average generation time, an improvement of up to 50%. In GAMA, several constants must be calculated for each different value of the shape parameter A; these constants are saved between calls so that they need not be recomputed. It will thus be more efficient to get several gamma variates with the same shape parameter before changing the A value, especially when A > 3.0 when the setup computations are extensive (see lines

174-246 of the program listing).

Note that the techniques used in GAMA and CAUCHY make use of so-called rejection methods so that the number of (or exponential or normal) deviates needed to generate a single output deviate is random. When normal or exponential deviates are required by GAMA from LLRANDOM a vector of 10 deviates is called for; since not all of these may be used at the time they are generated, the balance are saved for the next call to GAMA. Thus, reinitializing the seed IX to its original value will not in general result an exact repetition of the generated gamma sequence since the first few deviates will use the old normal exponential deviates from the previous sequence. To achieve an exact repetition, the generator must be forced to repeat the initialization computations for the desired A value; at this time any remaining variates from LLRANDOM discarded. An example of this might be

CALL OVPLOW

IX = 12345

CALL GAMA ( A, IX, G, 100 )

C REINITIALIZE GAMMA SEQUENCE

CALL GAMA ( 1.0, IX, G, 1 )

IX = 12345

CALL GAMA ( A, IX, G, 100 )

CALL GAMA ( A, IX, G, 100 )

END

CAUCHY requires 552 bytes and, as mentioned previously, is completely independent of any other subprograms. CAUCHY uses the LLRANDOM multiplicative congruential uniform

generator but this is coded in line when needed so as to preserve CAUCHY's independence. The average generation time per variate for subroutine CAUCHY on a System/360 Model 67 under OS/MVT was 67.5 microseconds when variates were generated in vectors of 100. The generation of variates one at a time increased the average time to 119.3 microseconds per variate.

Subroutine GAMA itself uses only 1988 bytes of memory but since it calls on LLRANDOM the total core requirement is 9342 bytes:

GAMA			1988	bytes
LLRANDOM			6189	bytes
Required	IBM	<b>Functions</b>	<u>1165</u>	bytes
Total			9342	bytes

Timing the gamma generator on a System/360 Model 67 was carried out using the TIME macro; Table 1 summarizes the observed times as a function of the shape parameter, A. Note that since special methods are employed when A is 0.5, 1.0, 1.5, 2.0 or 3.0, the times in these cases are considerably shorter than times for nearby values of A.

Shape	Parameter	Algorithm	Vector of 100	Single
	<b>A</b>		Variates ·	Variate
	0.1	GS	324.0	364.0
	0.3	GS	367.0	402.5
	0.5	GA	70.4	207.7
	0.8	GS	439.8	551.2
	0.9	GS	459.0	611.0
	1.0	G <b>A</b>	68.7	158.9
•	1.2	GF	300.1	385.0
	1.4	<b>GP</b>	306.1	441.0
	1.5	GA	141.7	215.8
	1.8	<b>GP</b>	343.6	390.8
	2.0	GA	142.5	203.6
	2.1	<b>GP</b>	396.1	450.8
	2.5	G <b>P</b>	434.7	468.5
	2.9	<b>GF</b>	444.5	. 496.6
200	3.0	GA	206.7	237.1
	3.1	GO	341.5	435.8
	3.5	GO	336.2	373.4
	4.0	GO	332.4	420.7
	5.0	GO	307.7	363.2
	8.0	GO	293.1	371.3
1	10.0	GO	289.4	312.5
2	20.0	GO	238.2	321.6
5	0.0	GO	197.7	284.2
10	0.0	GO	178.4	220.0
100	0.0	GO	166.7	177.0
1000	0.0	GO	136.4	169.8
10000	0.0	GO	152.5	235.8

Table 1. Average generation times (microseconds) for gamma variates using subroutine GAMA.

## III. Description of the Algorithms

This section describes the actual algorithms used in CAUCHY and GAMA. An understanding of the algorithms is not necessary for use of the package but they are set forth here both in the interest of completeness and in an effort to document the programs more fully. A single algorithm suffices for the Cauchy generator while GAMA uses one of four algorithms, depending on the value of A.

In the descriptions which follow, the letters U, N and E (with or without affixes) represent uniform, standard normal and unit exponential pseudorandom deviates, respectively. The phrase "Generate U" implies that U is the next sequential uniform variate in the linear congruential sequence; these variates are generated as needed by using the same multiplicative congruential scheme as used in LLRANDOM. The phrases "Generate N" or "Generate E" imply that normal or exponential variates are to be obtained by linking directly to LLRANDOM.

#### A. Cauchy Generator

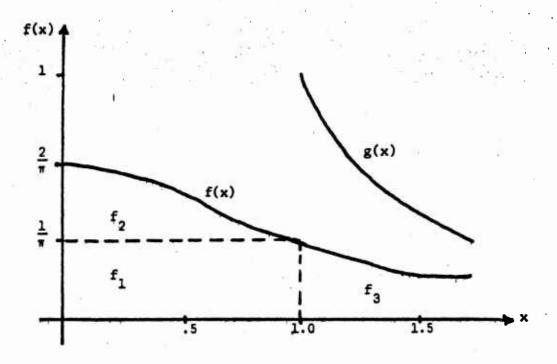
The Cauchy generator is a combination decomposition-rejection method (see Knuth [5]). The Cauchy density is decomposed, as in Figure 1, into three subdensities: a uniform density between 0 and 1  $(f_1)$ , a wedge-shaped density  $(f_2)$ , and a long tailed density  $(f_3)$ .

The uniform density f is sampled with probability  $1/\pi$ ; in this case a uniform (0,1) variate is returned. The density f is dealt with by using Marsaglia's almost-linear

density algorithm, just as in Knuth's Algorithm L [5]. The density  $f_2$  is sampled with probability  $1/2 - 1/\pi$ . The tail density  $f_3$  is sampled by a rejection method with probability 1/2. The majorizing density for  $f_3$  is  $g(x) = 1 / x^2$ , which is the density of the reciprocal of a uniform (0,1) variate.

Algorithm C below uses the fact that in the prime modulus congruential random number generator used in LLRANDOM the low order bits are uniformly distributed so that b and b select the proper sub-distribution in Step 1.

This will not in general be the case for other congruential pseudo-random number generators.



Pigure 1. Decomposition of the Cauchy Density Punction.

## Algorithm C. Cauchy variates.

- (Select subdensity) Generate U, setting aside the two low order bits b and b. If b = 1, go to Step 6.
- 2. (Sample box) If  $U \le 0.6366197724 = 2/\pi$ , generate a new variate U, set x = U and go to Step 8.
- 3. (Sample wedge) Generate new variates U and U. If U

  > U, exchange U and U. Set x = U.
- 4. (Easy rejection) If  $0 \le 0.8284271247 = 2\sqrt{2} 2$ , go to Step 8.
- 5. (Hard rejection) If  $U_2 U_1 \le \frac{1 x^2}{1 + x^2}$  ( 2/2 2), go to Step 8, otherwise go back to Step 3.
- 6. (Sample tail) Set x = 1 / U.
- 7. (Tail rejection) Generate a new variate U. If U  $\leq \frac{x^2}{1+x^2}$  go to Step 8, otherwise generate a new U and go back to Step 6.
- 8. (Random sign) If b = 1 set x = -x. Deliver x as the generated deviate.

It should be noted that there are several other methods for generating Cauchy variates: the ratio of independent standard normal deviates has the Cauchy distribution, as does the quantity

$$X = \tan \left[ \pi \left( \sigma - \frac{1}{2} \right) \right],$$

where U is uniform (0,1). These methods are both substantially slower than algorithm C, but another new method has an

average time comparable to Algorithm C and is much easier to program. This second method requires an average of 2.55 uniform random variates per Cauchy variate (as compared with 2.47 for algorithm C) and it needs about 69 microseconds per variate on the System/360 Model 67. It is possible, however, that Algorithm CR will be better than algorithm C in some other implementation.

The method is essentially the technique devised by von Neumann to generate a random variate  $\sin U$ , where U is uniform between 0 and  $2\pi$ . Such variates are used in the polar method for generating normal random variables [8]. It does not seem to have been recognized that the method also generates tan U, which is the required Cauchy variate.

# Algorithm CR. Cauchy variates, ratio method.

- 1. (Get uniforms) Generate U and U. Set Y = 2 U 1
  and Y = 2 U 1.
- 2. (Rejection test) If  $Y^2 + Y^2 > 1$  go back to Step 1.
- 3. (Take ratio) Deliver  $x = Y_1 / Y_2$ .

# B. Gamma Generator GS: A < 1.0

This method is due to Ahrens and is set forth in [1]. It is applicable only to values of A less than one and is markedly superior in execution time to the method of Johnk [3], which is the usual technique for generating variates of this type.

The method is a rejection method employing two different tests, one of which is chosen at random for any given variate: the power transform of a uniform (0,1)

variate,  $U^{1/\lambda}$ , is tested in the region 0 < x < 1, while a suitable exponential, E, is tested when x > 1. The advantage of this method lies in the limited use of the library subprograms for the exponential and logarithm; average times range from 300 to 400 microseconds as compared with 600 to 800 for Johnk's method. Further discussion and proofs may be found in [1].

# Algorithm GS. Gamma variates, A < 1.0.

- 1. (Select rejection test) Generate U and generate E and set P = e + A U. (Note that "e" is the base of the natural logarithms.) If P ≤ 1 go to Step 2, otherwise go to Step 3.
- 2. (Small x test) Set  $x = P^{1/A}$ . If  $x \le E$ , deliver x, otherwise go back to Step 1.
- 3. (Large x test) Set  $x = -\ln \left[ \frac{1}{A} \left\{ \frac{e+h}{e} P \right\} \right]$ . If (1 A)  $\ln x \le E$ , deliver x, otherwise go back to Step 1.

# C. Gamma Generator GF: 1.0 ≤ A ≤ 3.0

A thus-far unpublished method devised by Professor G.S. Fishman of North Carolina University was communicated to the authors in private correspondence. It is valid for any  $\lambda > 1.0$  but its efficiency in terms of average time goes down as  $\sqrt{\Lambda}$  so it is applied in GAMA only in the range where it is superior to the Dieter-Ahrens method GO described below.

The method is a rejection method based on the following theorem.

Theorem Let U be a uniform (0,1) random variable and let E be an exponential random variable with mean A. Let

$$g(x) = \begin{bmatrix} x \\ 1 \end{bmatrix} = e^{-x(1-1/\lambda)} - (\lambda-1)$$

If  $g(E) \ge U$ , then E has conditionally the gamma distribution with shape parameter  $\lambda$ , i.e.

$$f_{E}(x)U \leq g(B)) = \frac{A-1-x}{\Gamma(A)} -$$

Proof:

Unconditionally, E has density  $h(x) = \frac{1}{A} e^{-x/A}$ .

Therefore,

(4) 
$$f_{E}(x \mid U \leq g(E)) = h(x) Pr(U \leq g(E) \mid E = x)$$

Now since U is uniformly distributed,

$$Pr\{U \leq g(E) \mid E=x\} = g(x)$$

as long as 0 < g(x) < 1; that this is true for every x > 0 may be readily verified by elementary calculus. Therefore,

(5) 
$$\Pr\{U \leq g(E)\} = E[\Pr\{U \leq g(E) \mid E\}]$$

$$= \int_0^\infty g(x) h(x) dx^4$$

$$= \int_0^\infty f(x) dx dx^4$$

Thus, in view of (4),

$$f_{E}(x) \cup g(E) = h(x) \cdot g(x)$$

$$= \frac{A-1}{\Gamma(A)} - x$$

The efficiency of the generator is governed by the probability that a given variate will pass the rejection test,  $U \le g(E)$ ; from (5) it will be seen that this probability is just  $C(\lambda)$ . When  $\lambda$  is large we have from Stirling's approximation that  $C(\lambda) = \sqrt{\frac{2\pi}{\lambda}}$ , so that the method becomes more inefficient with increasing  $\lambda$ , as noted above.

A slight modification to the method suggested by the theorem improves the efficiency slightly and we obtain

Algorithm GF. Gamma variates, 1.0 < A < 3.0.

- 1. (Generate exponentials) Generate two independent exponential variates,  $E_1$  and  $E_2$ .
- 2. (Rejection test) If  $E_2 < (\lambda-1)$  ( $E_1 \ln E_1 1$ ) then go back to Step 1.
- 3. (Acceptance) Deliver  $x = \lambda E_1$ .

#### D. Gamma Generator GO: A ≥ 3.0

This method was originally developed by Dieter and Amrens and is fully described in [1] together with several other gamma generation techniques. Algorithm 60 does not

suffer the usual drawback of growing less efficient in generation time with increasing A; in fact, the method is more efficient for larger A values.

The basic idea here is to take advantage of the asymptotic normality of the gamma distribution by doing most of the sampling from a normal distribution; the right hand tail is sampled, when necessary, using a rejection method with the exponential distribution. The method can be applied to values of A greater than 2.533, but it is not as efficient as Fishman's technique for A < 3.0.

As mentioned previously, this algorithm requires the computation of several constants which depend only on A and which may be saved between calls; these calculations are described in step 0 of the specification below. Further discussion, illustrations and proofs are given in [1]; the version of GO here differs in a few minor details from the original Dieter and Ahrens technique.

# Algorithm GO. Gamma variates, a > 3.0.

O. (Calculate constants) Compute:

$$m = A - 1;$$
  
 $s^2 = \sqrt{8A} + A;$   $s = \sqrt{s^2};$   
 $d = \sqrt{6s^2};$   $b = d + n;$   
 $v = s^2 / m - 1;$   $v = 2s^2 / (m / A);$   
 $c = b + \ln \frac{s}{5} - \frac{d}{5} - 2m - 3.7203285.$ 

- (Select normal/exponential) Generate U. If U ≤
   0.0095722652 go to Step 7.
- 2. (Normal sampling) Generate N and set x = sN + m.
- 3. (Check trial value) If x < 0 or x > b go back to Step 2,

otherwise generate a new variate U and set  $S = N^2 / 2$ . If N > 0 go to Step 5.

- 4. (Left-hand rejection) If U < 1 + S (VN w) go to Step 9, otherwise go to Step 6.
- 5. (Right-hand rejection) If U < 1 ws go to Step 9.
- 6. (Final normal rejection) If  $\ln U < m \ln \frac{x}{m} + m x + S$  go to Step 9; otherwise go back to step 1.
- 7. (Exponential) Generate E and E and set x = b(1+E/d).
- 8. (Exponential rejection) If  $m (\frac{x}{b} \ln \frac{x}{m}) + c > E$  go back to Step 1.
- 9. (End) Deliver x as the gamma variate.

### E. Ad Hoc Gamma Generators

This set of algorithms is based on the well-known fact that the sum of independent gamma variates with shape parameters A and A and equal scale parameters has the 2 gamma distribution with shape parameter A + A and scale parameter equal to that of the summands. We may thus generate a gamma variate with integer shape parameter K by taking the sum of K independent exponentials. This will be more efficient than the previously discussed methods (Algorithms GF and GO) for moderate values of K; for the System/360 we take K ≤ 3 to apply this ad hoc technique.

An obvious extension to this method is to allow for half-integral values of A by making use of the fact that the square of a standard normal random variable has the chi-squared distribution with one degree of freedom, i.e.  $N^2/2$  has the gamma distribution with unit scale parameter and A = 0.5. We use this extension for A = 0.5 or 1.5.

# The resulting algorithm is then

- Algorithm GA. Gamma variates, integral or half-integral shape parameter A.
- 1. (Find K) Set K = [A], where [A] denotes the integral part of A. Set X = 0. If A K = 0.5 set L = 1; if A K = 0.0 set L = 0; otherwise Stop. (If the algorithm stops, an incorrect A value has been used.)
- 3. (Generate normal) If L = 0 go to Step 4 otherwise generate N and set  $X = X + N^2/2$ .
- 4. (Deliver X) X is the desired variate.

#### IV. Summary and Comments

This work provides a convenient and useful extension to the LLRANDOM package, especially for users interested in statistical and reliability theory applications of digital simulation. The combination of the most efficient known gamma generation techniques with the new Cauchy method gives exceptionally good time characteristics at some cost in computer memory utilization.

The work may be extended at once to the generation of several other types of random variables. For example, the beta distribution with parameters A and B may be sampled by taking gamma variates X and X with respective shape parameters A and B and delivering

$$z = x_1 / (x_1 + x_2)$$

as a beta variate. In this case considerable overhead in GAMA can result from shifting the shape parameter back and forth between A and B; for this reason obtaining vectors of gamma variates X and X is recommended, as in the following example:

DIMENSION X1(50), X2(50), Z(50)

...

CALL GAMA ( A, IX, X1, 50 )

CALL GAMA ( B, IX, X2, 50 )

DO 405 I = 1,50

Z(I) = X1(I) / ( X1(I) + X2(I) )

405 CONTINUE

END

The <u>t-Distribution</u> may be sampled as the ratio of a standard normal and an independent chi-squared random variate, while the <u>F-Distribution</u> may be obtained by taking the ratio of two independent chi-squared variates divided by their respective degrees of freedom. (See pages 4 and 5 for an example of the generation of chi-squared variates.)

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GE USAGE:	GENERALTION OF RANDOM VARIATES WITH I	
ŲSAGE		THE CAUCHY DISTRIBUTION
•	: 30	
<b>ی</b>	CALL CAUCHY (IX, C, N)	
PARAM	PARAMETERS:	
×	SEED FOR RANDOM NUMBER GENERATOR (INITIALIZED TO ANY POSITIVE VALUE AND NOT ALTERED THEREAFTER.	INTEGER#4). SHOULD BE IN THE CALLING PROGRAM
ِن	ARRAY TO HOLD THE GENERATED VARIATES DIMENSIONED AT LEAST N.	S (REAL*4). MUST BE
z	NUMBER OF CAUCHY DEVIATES TO GENERATE	NTE (INTEGER*4).
METHOD:	нор :	
SUBDI	A COMBINED DECOMPOSITION/REJECTION MISUBDISTRIBUTIONS CAN BE SAMPLED USING UN	HETHOD IS USED. ALL
SUBRO	SUBROUTINES REQUIRED:	
Z	NONE	
PROGR	PROGRAMMER: D.W. ROBINSON	
DATE:	E: 9 MAY 1974	

AU0038 AU0038 AU0038 AU0040 AU0041	AU00049 AU00049 AU000469 AU0046	AACOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOO	AU0053 AU0053 AU0054	AU0056	AU00058 AU00058 AU00058	AU0061 AU0062	AU0063 AU0064	AU0068 AU0067 AU0068	CCACCOSO CCACCOSO CCACCOSO CCACCOSO CCACCOSO CCACCOSO CCACCOSO CCACCOSO CCACCOSO CCACC
EGISTER ALLUCA 0 SAVE +/- 1 WORK REGI	R2 CONSIANT 4 R3 NUMBER OF DEVIATES (BYTES) R4 BASE ADDRESS OF C ARRAY R5 INDEX OF CURRENT RANDOM NUMBER IN C	R6,R7 SEED FOR GENERATOR R8 UNIFORM MULTIPLIER = 16807 R9 EXPONENT CONSTANT = 40000001 R10 NORMALIZATICN COMPARAND = 40100000	RII CONSTANT I (MASK) RIZ ADDRESS OF END OF MAIN LOOP	R13 ADDRESS OF IX IN CALLING PROGRAM	RI4 RETURN ADDRESS RI5 BASE REGISTER	**************************	UNIFORM RANDOM NUMBER GENERATION MACRO	WITH THE CURRENT UNIFORM INTEGER IN R7 AND THE MULTIPLIE IN R8, FINDS THE NEXT UNIFORM INTEGER AND PUTS IT INTO R	MACRO RAND RAND RAND RAND RAND RAND RAND RAND

\*\*\* CAUCHY DEVIATE GENERATOR \*\*\*

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AUCHY	ODBOOM NEC TOO M	CAUCHY, R15 12(, R15) AL1(6) CL6'CAUCHY' R14, R12, 12(R13)	DEFINE BASE REGISTER RANCH AROUND ID MODULE NAME SAVE CALLING PROGRAM REGS	AAAAA 000000 0000000000000000000000000	
	STAR-	13,5VAKEA+ 2,R13 13,5VAREA 13,8(,R2)	AVE ADDRESS IN D ING SAVE ADDRESS AREA IN R13 INK	AU0092 AU00992 AU009943 AU009943	
_	NI N	R3, R5, O(R1) R13, R3 R7, O(, R3) R3, O(, R5) R3, 2 R2, 4	ETER ADDRESSES ADDRESS VALUE ER DEVIATES TO BYTES 4 FOR MAIN LOO	CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC	
	CLER CNO	400 -	ACK UP 4 IN CALLER'S NITIAL ARRAY INDEX LOAD MAIN LOOP CONST LIGN BXLE LOOP FOR S	AU0105 AU0105 AU0105 AU0105	
4AINLGOP	RAND LR	-0-	GET FIRST UNIFORM SAVE TWO BITS OF X(N) LAST BIT OF X(N) IN RO NEXT TO LAST BIT IN RI	AAUOI 098	
	SRL NR BZ	RIFI RIFRII TAIL	EST BIT IN RI; IF O,	ACOLLIANT OF THE PROPERTY OF T	
	SE E	R6,=F'136713055 WEDGE	1 SELECT RECTANGLE/WEDGE SAMPLING	AU0116 AU0117 AU0118	
SAMPL	RAND SRL OR ST	64.7 67.7 6.0 6.0 6.0 6.0 6.0 6.0 6.0 6.0 6.0 6.0	GET NEXT UNIFORM MAKE ROOM FOR EXPONENT "OR" ON THE EXPONENT STORE THE UNIFORM	A A C C C C C C C C C C C C C C C C C C	
		ONW	TEST FOR NORMALIZATION QUIT IF NOT NEEDED NORMALIZE THE UNIFORM GO TO END OF LOOP	A W O O O O O O O O O O O O O O O O O O	

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AU0130 AU0131	AU0132 AU0133	AU0135	AAACOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOO	ACO14100141001410014100141001410014100141	AU0144 AU0145	CANCOLOR (CANCOLOR (CANCOL	AUOLUZA AUOLUZA AUOLUZA AUOLUZA	AU0156 AU0156 AU0156	AU0160 AU0160
	FORM IN RI		TO REAL	TO REAL		REJECTION TEST		SQRT(2) )	
SAVE FIRST UNIFORM	ET UNIFORM IN R6 < UNIFOR	XCHANGE REGISTERS	CCEPT WEDGE SAMPLE ONVERT MINIMUM UNIFORM OR THE EXPONENT	ONVERT MAXIMUM UNIFORM OR # ON THE EXPONENT	OAD TRIAL VARIATE EST FOR NORMALIZATION	ORMALIZE X ET FIRST COMPARAND FOR U2 - X IND X ** 2	X + 2 IN X + 2 X + 4 X + 2 X +	CON SNO SNO SNO SNO SNO SNO SNO SNO SNO S	D BACK IF TEST FAILED
ND RI,R6	R6,R1 G	R6,R1	RI = F 1779033703 SAMPL R6.7 R6.7 R6.19	R1,7	FRO UNIF	R FR2 FR0 6 6 6 FR2 FR0	ER FR6, FR4 FR6, ER 1.00	R FR6,FR4 FR6,=E1.82842712 R FR5,FR6 13.21,7	MEDG
270	202		್ಷ ಬ್ಲಿಕ್ಟ್ ಬ್ಲಿಕ್ಟ್			review Serview			

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MAKE ROOM FOR EXPONENT "OR" ON THE EXPONENT STORE THE UNIFORM GET 1 / UNIFORM FOR REJECTION TEST MAKE ROOM FOR EXPONENT	GET 1 + X ** 2  GET 1 + X ** 2  FIND COMPARAND FOR REJECTION TEST  REJECTION TEST  ANOTHER UNIFORM FOR NEXT PASS  GO BACK	TEST SAVED BIT  IF BIT = 0, QUIT  IF BIT = 1, X = -X  STORE VARIATE IN CALLER'S ARRAY  BRANCH BACK FOR ANCTHER VARIATE	SEND LAST SEED BACK TO CALLING PROGRAM GET CALLING SAVE AREA ADDRESS PRESTORE CALLING PROGREGS RETURN
R6.7 R6.89 R6.UNIF FRO.E!1.0 FRO.UNIF R6.7 R6.89	FR2, FR0 FR4, FR2 FR4, UNI F FR4, UNI F FR4, FR2 I3, R12	RO, R11 *+6 FRO, FRO FRO, O(R4, R5) RS, R2, MAINLOOP	R7,0(,R13) R13,SVAREA+4 R14,R12,12(R13 R14
SRL ST SAND SRL ST ST ST ST ST ST ST ST ST ST ST ST ST	7274200000 MMMMMMMM AXX C C	NR BZ STE BXLE	ST RE
AIL		ENDLOOP	<u>.</u>

	DATA	AREA		
SVAREA	DS	18F	SAVE AREA	
UNIF U2	DS DS	u.u.	TEMP STORAGE FOR UNIFORM RANDOM VARIATES	
LOGPCON	00 <b>000</b>	F'16807' X'40000001' X'40100000' F'1' AL4(ENDLOOP)	MULTIPLIER FOR GENERATOR EXPONENT CONSTANT NORMALIZATION TEST CONSTANT MASK CONSTANT END OF LOOP ADORESS	R8 R89 R10 R12 R12
	LTORG	G		
	<b>REGISTER</b>	STER EQUATES		·
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28888888		145-88-46-1 101-10		
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\*\*\*\* GAMMA DEVIATE GENERATOR \*\*\*

PUR POSE:

GENERATION OF PSEUDO-RANDOM GAMMA DEVIATES WITH NON-INTEGRAL SHAPE PARAMETER A > 0 AND SCALE PARAMETER I.

USAGE:

CALL GAMA (A, IX, G, N)

PARAMETERS:

GAMMA SHAPE PARAMETER (REAL#4). MUST BE >

SEED FOR GENERATOR (INTEGER\*4). SHOULD BE INITIALIZED IN THE CALLING PROGRAM TO ANY POSITIVE VALUE AND NOT ALTERED THEREAFTER. XI

ARRAY TO HOLD THE GENERATED DEVIATES (REAL\*4). SHOULD BE DIMENSIONED AT LEAST N.

DELIVERED (INTEGER\*4). NUMBER OF GAMMA DEVIATES TO BE

METHOD:

THREE DIFFERENT BASIC METHODS ARE USED, DEPENDING ON THE VALUE OF A:

AHRENS SMALL PARAMETER METHOD (ALGORITHM "GS") FISHMAN'S REJECTION METHOD (ALGORITHM "GF").

DIETER-AHRENS NORMAL-EXPONENTIAL METHOD (ALGORITHM "GO").

WHEN A IS EXACTLY 0.5, 1.0, 1.5, 2.0 OR 3.0 AN AD HOC METHOD BASED ON TAKING THE SUM OF INDEPENDENT EXPONENTIALS IS USED.

\*\*\*\* GAMMA DEVIATE GENERATOR \*\*\*

SUBROUTINES REQUIRED:

PACKAGE ALOG, GENERATOR FUNCT I ONS LEWIS AND LEARMONTH RANDOM NUMBER 1 IS NEEDED. THE FORTRAN BUILT-IN SORT ARE ALSO USED. LLRANDOM I EXP AND SQ

NOTES:

EO BECAUSE THE GENERATED DEVIATES WILL BE TOO SMALL. THE FORTRAN STANDARD FIXUP IN THIS CASE IS TO SET THE GENERATION OF THIS CASE IS TO SET THE GENERATION OF THIS MAY CAUSE PROBLEMS IF FURTHER DATA TRANSFORMATIONS (E.G., LOGARITHMS) ARE PLANNED.

LARGE ⋖ IF **EFFICIENT** 2. THIS SUBROUTINE IS, IN GENERAL, MORE NUMBER OF GAMMA DEVIATES IS GENERATED.

NORMAL OR EXPONENTIAL DEVIATES BY METHODS GO, GS, OR GF, IT MAY TWO COMPLETELY DIFFERENT SEQUENCES SEEDS. MILL BE SAVED BETWEEN CALLS NOT BE POSSIBLE TO PRODUCE OF DEVIATES WITH DIFFERENT S

PROGRAMMÉR: D.W. ROBINSON

27 JANUARY 1975

**VERSION** 

1 ADDED 0.5, 1.5, 2.0 AND 3.0 METHODS

\*\*\*\* GAMMA DEVIATE GENERATOR \*\*\*

GENERATO	EQUATES:	
VIATE	いることであることでは、これでは、これでは、これでは、これでは、これでは、これでは、これでは、これ	074
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\*\*\*\* GAMMA DEVIATE GENERATOR \*\*\*

LINKAGE / INITIALIZATION SECTION

GAMA

MODULE IDENTIFIER
SAVE CALLING REGS
CALLING SAVE ADDRESS IN OWN AREA
COPY CALLING AREA ADDRESS TO RZ
OWN SAVE AREA IN R13 DEFINE BRANCH GAMA, R15 10(1815) AL1(4) CL4.(GAMA, R14, R12, 12(R13) R13, SVAREA+4 CC R2, R13 R13, SVAREA 

BASE REGISTER AROUND ID

R2, R5, O(R1) FRO, AP SETUP STUP R7, O(, R3) R3, O(, R5) R3, O(, R5) R4, R2 R5, R2 R6, METHOD

GET PARAMETER ADORESSES
GET SHAPE PARAMETER
TEST FOR NEW "A" VALUE
IF SO, DO PRELIMINARY CALCULATIONS
CONSTANT 4 FOR MAIN LOOP
PUT SEED INTO R7
GET NUMBER OF DEVIATES, N
CONVERT TO BYTES
BACKUP ONE IN CALLER'S ARRAY
INITIAL MAIN LOOP INDEX
JUMP TO PROPER METHOD

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SETUP

\$		OOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOO
		·
AND CONSTANT CALCULATION  FRO, FRO THRU FRO, AP FRO, A	P FOR LARGE PARAMETER METHOD, ALGORITHM "GO"  RO, GO  RO, GO  RO, HOD  RO, HOD  INITIALIZE RANDOM ARRAY INDEX  FRO, AGJ  GJ AHEAD IF NOT  FRO, AGO  SAVE NEW SHAPE PARAMETER  SAVE NEW SHAPE PARAMETER  COMPUTE MU = A - 1.  FRO, FRZ  FRO, MU  FRZ, FRO  COMPUTE MUP = 1 / MU	RI, ARGLSTI LOAD ARGUMENT LIST R8, R15 R15, VADDSR ADORESS OF SQRT FUNCTION R15, VADDSR ABORESS OF SQRT FUNCTION R15, R8 R15, R8 R15, R8 SAVE SQRT (A) FR2, FR0 FR0, = 1.6329932* FIND NORMAL VARIANCE FR0, SIGMA
2	SON STEER THE STANK OF THE STAN	NAMERIAN TO THE STATE OF THE ST

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TTTTTTTTT	CXXXXXXXXX	AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA	EXXXXXXXX	EXXXXXXXX
TANT "WM" CONSTANT "VP"	ADDRESS CTION ER CONSTANT "DP"	NORMAL METHOD, "8"	CONSTANT "CONS"  OR LOG FUNCTION  ESS  NCTION  SS	N OF "CONS" ATION. PROCEED TO
FIND REJECTION CONS 2. FIND REJECTION TO FIND NORMAL STO	LOAD ARGUMENT LIST ADDRESS OF SQRT FUN RESTORE BASE REGIST SAVE STD DEV	FIND UPPER LIMIT FOR COMPUTE BP = 1 / B	COMPUTE REJECTION COR FIRST FIND VALUE FOR LOAD ARG LIST ADDRESS ADDRESS OF ALOG FUNC RESTORE BASE ADDRESS	OMPLETE COMPUTATIO  ONE WITH INITIALIZ  GENERATION
FRO, MU FRO, MM FRZ, MM FRZ, MU FRZ, MU FRZ, VP TO SORT FUNCTION	RRISTAND CAN THE CONTROL OF THE CONT	MU BE BE BE	FR2.SIGMA FR2.D FR2.FR0 FR2.CONS RI.ARGLST3 RI5.VADDLG RI5.R8	FRO, FRO FRO, B FRO, MU FRO, EE 3. 7203285 FRO, CONS GWAN
M M M M M M M M M M M M M M M M M M M	• <b>4 44 4 1 1 1 1</b>	SOFE SOFE SOFE SOFE SOFE SOFE SOFE SOFE	ME DOER STE LA LA BALR	BNAAMER TAMERER BNAAME

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A ROPE	E FROTAMINUS INITIALIZE RANDOM ROTON SON INX2 DONE WITH INITIAL GWAN GENERATION.	METHOD.  DDRESS FO  FE 1 - A  FE 1 / A  IND (E +	TO GENERATION.

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SET	UP FOR AD HOC MET	METHODS	666 666 668 668 668 668 668 668 668 668
SET	UP FOR CHI-SQUARED,	ED, 1 DEGREE OF FREEDOM ( A = 0.5 )	MA 28
LA	RO, CHISO1	SET ADDRESS FOR SUBSEQUENT CALLS	A A A A A A A A A A A A A A A A A A A
	MAN	GO ON TO GENERATION	MA 28
SET	UP FOR EXPONENTIAL	AL (A = 1.0)	MA 29
Y L		SET ADDRESS FOR SUBSEQUENT CALLS	MA 29
- 0 0	NA N	GO ON TO GENERATION	MA 29
SET	UP FOR CHI-SQUARED,	ED, 3 DEGREES OF FREEDOM ( A = 1.5 )	MA 29
	O	SET ADDRESS FOR SUBSEQUENT CALLS	MA 29
- 4t	E <	INITIALIZE RANDOM ARRAY INDEX	MA W
	NAN NA	GO ON TO GENERATION	A A A
SET	UP FOR 2 - ERLANG	G (A = 2.0)	A A A
	O C	SET ADDRESS FOR SUBSEQUENT CALLS	AA WW
- <b>4</b> h	000 000 000 000 000 000 000 000 000 00	INITIALIZE RANDOM ARRAY INDEX	200 7 M M
	MAN	GO ON TO GENERATION	MA 31
SET	UP FOR 3 - ERLANG	G ( A = 3.0 )	MA A
	O.	SAVE ADDRESS FOR SUBSEQUENT CALLS	MA AN
- Al	C u	INITIALIZE RANDOM ARRAY INDEX	AAA VOC
	MAN	GO ON TO GENERATION	MA 31

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OQOQQQQQQQQQQQQQQQQQQQQQQQQQQQQQQQQQQQ	ETTTTTTTTTTTTTT	PETETETETE
METHOD "GO" (DIETER-AHRENS)  LM R8*R13, GOCON LOAD LOOPING CONSTANTS  CNOP 0,8  MR R6.R8  SLDA R6,1  SRL R7,1  AR R6,R7  BNO R4+10  BNO R4+10  AR R6,R2  LR R7,R5  CO ON IF NO OVERFLOW  AR R6,R2  LR R7,R6  CO ON IF NO OVERFLOW  AR R6,R2  LR R7,R6  CO ON IF NO OVERFLOW  AR R6,R2  BUT X(N) INTO R7  SAMPLING FROM THE NORMAL DISTRIBUTION	BXLE R12,R1G,GONTST INCREMENT NORMAL ARRAY INDEX.  ST R7,1X  SAVE CURRENT SEED VALUE.  LA R13,8VAREA SAVE BASE REGISTER  LA R15,VADDNM ADDRESS OF NORMAL GENERATOR  LA R15,VADDNM LINT TO "NORMAL GENERATOR  LA R15,VADDNM CONTRACT ARGURENT TO "NORMAL GENERATOR  LA R15,VADDNM CONTRACT ARGURENT TO "NORMAL GENERATOR  LA R13,ENDGO RESTORE END OF LOOP REGISTER  SR R12,R12  RESTORE END OF LOOP REGISTER  R7,1X  RESTORE SEED  ALIGN BXLE LOOP FOR SPEED	LER FRO, RNARRAY(R12) LOAD NEXT NORMAL DEVIATE LER FR2, FRO ME FR2, SIGMA X = NORMAL * SIGMA + MU AE FR2, MU AE FR2, MU BNP GONORM CE FR2, B GONORM REJECT X < 0 ER4, FR0 S2 = 0.5 * S * S HER FR4, FR0
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\*\*\*\* GAMMA DEVIATE GENERATOR \*\*\*\*

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	EXXXXXXX	EXXXXX	EXXXXXX	EE	EXXXXXXX
NORMAL REJECTION TEST GET NEXT UNIFORM R6 = REMAINDER; R7 = QUOTIENT ADD QUOTIENT TO REMAINDER THUS SIMULATING DIVISION BY 2 ** 31 - 1 GO ON IF NO OVERFLOW. ADD 2 ** 31 - 3 ADD 4 MORE POLY (NO RY MAKE ROOM FOR EXPONENT. SAVE THE UNIFORM. PERFORM THE EXPONENT ON THE EXPONENT SAVE THE UNIFORM.	COMPUTE THE REJECTION VALUE:  1 + S2 * (S * VP - WM)  REJECTION TEST GO TO LOOP END IF PASSED. FURTHER TEST IF NOT.	COMPUTE THE REJECTION VALUE:  1 - S2 * WM  REJECTION TEST GO TO LOOP END IF PASSED.	FIND PARTIAL SUM FOR REJECTION TEST:  SUM = MU - X + S2  SAVE TRIAL GAMMA DEVIATE  GET LOG ARGUMENT, X / MU	TINE TWICE	SAVE PROGRAM REGS SAVE BASE REGISTER SAVE AREA POINTER ARGUMENT LIST ADDRESS ADDRESS OF FORTKAN LOG FUNCTION RESTORE BASE REGISTER
IFORM FOR 11	FR0,VP FR0,FR4 FR0,=E°1.0° FR0,UNIF 2,R13	FRO, FR4 FRO, WM FRO, UNIF 2, RI3	FR4, FR2 FR4, SUM FR2, X FR2, X FR2, TUP	TO LOG SUBROUTIN	RI2, RI3, GOSAV RI3, SVAREA RI3, VADDLG RI5, VADDLG RI5, RI5
MAR SSLOA SSROA BAR SSRL SSRL BP ER		CHIMING	SAN	LIR	LA L
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ADD MU * LOG (X / MU) TO SUM GET REJECTION VALUE SECOND LINK TO LOG FUNCTION ADDRESS OF LOG FUNCTION	ESTORE BASE REGIST RESTORE OTHER REG	RELOAD TRIAL GAMMA FINAL REJECTION TEST PASSED TEST. GJ TO LOOP END. FAILED TEST. BRANCH BACK FOR ANOTHER TRY.	FROM THE EXPONENTIAL DISTRIBUTION.		SAVE SECU. SAVE BROGRAM REGS. SAVE BASE REGISTER. SAVE AREA POINTER	ORESS OF EXPONING TO "EXPON" STORE BASE REG	FIND TRIAL GAMMA VALUE: X = B * (1 + R * DP)	SAVE TRIAL GAMMA VALUE GET LOG (X / MU)	ARGUMENT LIST SS OF LOG FUN TO "ALOG"	RESTORE BAS RESTORE OT
RO MU RO SUU 1 ARG	R15 R12 R13,6	FR2, X FR0, SUM 13, R13 601, 00P	REJECTION SAMPLING F	R7, IX	R12, R13, GOSAVE R12, R15 R13, SVAREA R13, SVAREA	15, VADDE 14, R 15 15, R 12	FRO, RNEXP FRO, DP FRO, = E' 1.0'	2000	450-5	12
SAPE TA	BAL LRAL	<b>1088</b> 円円 R	REJ	OEXP ST	IL A L		TAA	S X X	LA L	25
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\*\*\*\* GAMMA DEVIATE GENERATOR \*\*\*

GMA 4540 GMA 4550	NA 45	AA 450	MA 46 MA 46	MA 46	NA 46	MA 46	MA 46	MA 46	MA 40
RELOAD TRIAL GAMMA VALUE COMPLETE CALCULATION OF REJECTION VALUE. MU * (10G - x * RP) * CONS			TES	ACK TO START IF FAI	L00P.	e to in the	DEVIATE IN CALLER'S	KANCH BACK FUK ANDIHEK DEVIAL AVE 1 AST ABBAY INDEX	LL DONE DUIT
LE FR2,X LER FR4,FR2 ME FR4.8P	TRO TR		TRO TRO	00 TOO	END OF METHOD "GO" L	GENERAL ED DEVIA	FR2	ALE ROPRZIOLUU T R12.INX1	THRU
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XXX1	AAAAAAAA CARRESTA	II	EXXXXXXX	EXXX	EXXX	ZZZZ	INTERES.	EXXXX
	SET UP SEED LOAD LOOP CONSTANTS SHIFT BASE REGISTER KEEP "ALOG" ADDRESS IN RIS	ET NEXT PAIR DE EXPONENTIALS	ADDRESS GENERAT S TO R15 SPEED	12) TAKE LOGA DEVIATE	LINK TO "ALUG" RI2) FINISH COMPUTING REJEC (A - 1) * (R - LN R -	20(R12	DELIVER A * R STORE DEVIATE IN CALLER'S ARRAY BRANCH BACK FOR ANOTHER DEVIATE RESTORE BASE REGISTER	RELOAD SEED SAVE LAST ARRAY INDEX QUIT
ISHMAN'S METHOD	R7.IX R8.R12, GFCON R7.R15 R15 G GAMA,R7 R15,R9	R12, R10, GFTS	RI5, ARGL ST4 RI5, RB RI4, RI5 RI5, R9 R12, R12	6, RNARRAY (R 6, GFLOG	RI44RIS FRZ9RNARRAY( FR49FRZ FR49FRQ	R4, = E 1 . 0 . R4, AMINUS R4, RNA RRAY+ FLOOP	RARA SARA SARA SARA	G TRACT
FISH	CLUDER CRUDE	BXLE	LA LA SR CNOP		LEAL FER FR	TWWW	E X X 4 E	
***	F	GFL 00P		GFTST			*	

\*\*\* GAMMA DEVIATE GENERATOR \*\*\*

TETTETT	EXXXXXXXXX	EXXXXXXXXX	TATATATATATATATATATATATATATATATATATATA
2.0 OR 3.0 GREE OF FREEDOM ( A = 0.5 )	SAVE BASE REGISTER SKIP OVER SHAPE PARAMETER IN ARG LIST LINK TO "NORMAL" RESTORE BASE REGISTER GET SEED VALUE IN REG 7 ALIGN BXLE LOOP FCR SPEED	GET NEXT NORMAL SQUARE THE NORMAL AND MULTIPLY BY 0.5 PUT GAMMA DEVIATE INTO CALLER'S ARRAY BRANCH BACK FOR NEXT NORMAL QUIT	AVE B KIP D INK D INK D ESTOR ET SE
AD HOC METHODS A = 0.5, 1.0, 1.5, 2. CHI - SQUARED, 1 DEGR	LR R12,R15 LA R15,VADONM BALR R15,VADONM LR R15,R15 L R7,0(,R1) CNDP 0,8	LE FRO, 0 (R4, R5) HER FRO, FRO SIE FRO, 0 (R4, R5) BXLE R5, R2, CHLOOP1 EXPONENTIAL METHOD	RRIS 44 RRIS 4
****	CHISQ1	CHL 00P 1	<i>Z</i> A

\*\*\* GAMMA DEVIATE GENERATOR \*\*\*

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\*\*\*\* GAMMA DEVIATE GENERATOR \*\*\*

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II	EXXXX	E I	EXIX	EX	EXE:	EE	EXXXX	I	OCC SEE SEE SEE SEE SEE SEE SEE SEE SEE S
	SHIFT BASE REGISTER SKIP OVER SHAPE PARAMETER IN ARG LIST LINK TO "EXPON"	GET LAST SEED VALUE USED	SAVE SEED VALUE LOAD LOOP CONSTANTS ALIGN BXLE LOOP FOR SPEED	GET NEXT EXPONENTIAL	NEN LAL PARKAT EAR TO MEXPON!! ARGUMENT LIST	ESET ARRAY IN	J LOAD NEW EXPONENTIAL ADD TO SECOND EXPONENTIAL STORE GENERATED GAMMA IN CALLER'S ARRAY CO AACK EOB NEXT DEVIATE	O CHOCK I ON MEN I DEVINE	LOAD LAST SEED VALUE SAVE RANDOM ARRAY INDEX RESTORE BASE REGISTER
ERLANG ( A = 2.0	R6,R15 R1,4(,R1) R15,VADDEX	R7.0(,	7, IX 10, R12	E RIZ, RIO, CH4COMP	RIS, VADDEX RI, ARGLST4	12, RI	FRO, RNARRAY (R12 FRO, O(R4, R5) FRO, O(R4, R5)	Notice of the second	R7, IX R12, INX4 R15, R6 THRU
2 -	PLER.	Ž L	CNO	BXL		4 0K	SATE ATE		BEST
***	CHI SQ4		ŧ	CHL00P4	÷	*	CH4COMP	*	

\*\*\* GAMMA DEVIATE GENERATOR \*\*\*

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II	OOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOOO	EXXX	EZZ	EXXXX	EZZZ	EFE	EXXX:	EXX
			<b>+</b>					
	WASE REGISTER VER SHAPE PARAMETER IN ARG LIST O "EXPON" ST SEED VALUE USED	EED VALUE LOOP CONSTANTS BXLE LOOP FOR SPEED	EXT PAI	#ACK DAN # 1 EATAGS ED & ACK LENT MEKPAN LIST #EXPON" RAY INDEX	AD NEW EXPONENTIAL ADD TWO INDEPENDENT EXPONENTIALS	ENERATED GAMMA IN CALLER'S ARRAY K FOR NEXT DEVIATE	AST SEED VALUE ANDOM ARRAY INDEX E BASE REGISTER	
= 3.0 )	SHIFT SKIP O X LINK T GET LA	CHICON6 LOAD	GET	X LINK T 4 GET AR LINK AR RESET	AA	R5) SAVE G 00P6 GO BAC	LOAD L SAVE R RESTOR	1100
ERLANG ( A	R6.R15 R15.VADDE R14.R15	7, IX 10, R12,	R12, R10, C	RIS, VADDE RI, ARGLST RI4, RI5 RI2, RI2	O, RNAR	5, R2, C	R7, IX R12, INX5 R15, R6	G GAMA, R15 THRU
3 - 6	LA LA BALR	ST CNOP	BXLE	LA BALR SR			ST LR PB PB PB PB PB PB PB PB PB PB PB PB PB	SIN
***	CH1 SQ6		CHLOOP6		CHECOMP	4	·	

\*\*\*\* GAMMA DEVIATE GENERATOR \*\*\*\*

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	IIIIIII	EXXXXXX	EXIXX	EXX	EXEX.
MGS" (AHRENS)  OAD LOOP CONSTANTS LIGN BXLE LOOP FOR SPEED  ET NEXT UNIFORM DEVIATE 6 = REMAINDER; R7 = QUOTIENT 6 = REMAINDER; R7 = REMAINDER 6 = REMAINDER; R1 = REMAINDER 6 = R	LOAD FUNCTION ADDRESSES SHIFT BASE REGISTER TO R6	GET NEXT EXPONENTIAL IN ARRAY EXPONENTIAL ARRAY EXHAUSTED. REPLENISH IT SAVE SEED VALUE LOAD ARGUMENT LIST ADDRESS LINK TO "EXPON"	RESET ARRAY INDEX TO START RELOAD P INTO FRO RESTORE SEED TO R7 ALIGN BXLE FOR SPEED	FIND REJECTION METHOD TO USE	FIND LOG (P). LOAD ARGUMENT LIST ADD ADDRESS OF LOG FUNCTION
PARAMETER METHOD R8,R12,GSCON L 0,8 R6,1 R6,1 R6,1 R6,1 R6,1 R6,1 R6,1 R6,R7 R6,R7 R6,R9 R6,R9 R6,R9 R6,R9 R6,R9 R6,R9 R6,R9	R8, R9, GSVCON R15 GAMA, R6 E FROM EXPONENT	R12,R10, R7,IX R1,ARGLS R15,VADD	S A X	FRO,=E'1.0' XBIG	RI, ARGL ST9 RI5, R9 RI4, RI5
SMALL CNOP SRLDA SRL AAR ORL SRL SRL STE	LM LR DROP US ING SAMPLE	$\mathbf{c} \times \mathbf{r} \blacktriangleleft \blacktriangleleft$	CNOP	BE BH	LA LR BALZ
65 65 65 68 68 68	. **	* *		GSTST	хго

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	TITITITITITI		EXXXXXXXXXXXX	EXXXXXXXXX
			VALUE	
***	GET LOG (P) / A LINK TO EXPONENTIAL FUNCTION. LOAD ARGUMENT LIST ADDRESS RESULT IS P ** (1 / A) 2) REJECTION TEST QUIT IF OK; OTHERWISE GO BACK RESET BASE REGISTER	FIND (B - P) / A  NOW LINK TO LOG FUNCTION: ADDRESS OF LOG FUNCTION RESULT IS LOG ( (B - P) / A ) TRIAL GAMMA IS - LOG NOW FIND LOG OF TRIAL VALUE LOAD ARGUMENT LIST ADDRESS ADDRESS OF LOG FUNCTION	FINISH CALCULATION OF REJECTION 2) REJECTION TEST RELOAD TRIAL GAMMA VALUE QUIT IF OK OTHERWISE RESET LOOP CONSTANTS AND CHANGE BASE REGISTER AND GO BACK	IS IN FRO STORE DEVIATE IN CALLER'S ARRAY RESET LOOP CONSTANIS SHIFT BASE REGISTER BRANCH BACK FOR ANOTHER DEVIATE SAVE LAST ARRAY INDEX OTHERWISE QUIT.
IATE GENERATOR	FRO, AINV RIS, RB RI, ARGL ST9 RI4, RI5 FRO, RNARRAY(RI ENDGS RB, R9, GSCON RI5, R6 GSLOOP		SARMENTAR SARGEDOO SARGEDOO	VARIATE VALUE RA, RO, O(R4, R5) RB, R9, GSCON RIS, R2, GSLOOP RIZ, INX3 THRU R6
GAMMA DEVI	BRAPA CBAPA BRAPA	PLLSTORE BY SERVING BY		S SAMMA STER STER STER STER STER DE USROP
* * *		XBIG	***	ENDGS

\*\*\*\* GAMMA DEVIATE GENERATOR \*\*\*

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WEEKE STREET	EXX:	EXX	EXXXX:	EXXX:	EX:	EXXXXXXXXXX	
RESTORE CALLING SAVE AREA. GET ARGUMENT LIST ADDRESS GET SEED ADDRESS SEND BACK LAST SEED USED. ) RESTORE CALLING REGS	SAVE AREA	OLD SHAPE PARAMETER ADORESS FOR PROPER METHOD	EXTERNAL EXPONENTIAL GENERATOR EXTERNAL NORMAL GENERATOR LOGARITHM FUNCTION SQUARE ROOT FUNCTION	RANDOM NUMBER SEED ARRAY FOR NORMAL OR EXPONENTIAL DEVIATES NUMBER OF DEVIATES TO BE DELIVERED	#60#	SHAPE PARAMETER NORMAL MEAN NORMAL STD DEV UPPER LIMIT FOR NORMAL I / MU I / B MISC MISC CONSTANTS "GO"	
END OF ROUTINE.    R13,5VAREA+4	DS 18F	DS E'-1.0'	DC V(EXPON) DC V(NORMAL) DC V(ALOG) DC V(SQRT)	DS FOR TOP	CONSTANTS FOR METHOD	DC E 2.9413405 DC E 11.204783 DC E 0.25 DC E 0.89247598 DC E 1.9345306 DC E 1.9345306	
H C	SVAREA	AP METHOD	VAD PEX VADE CEX VADES CE	IX RNARRAY NUM	* * +	CVEDBMBUTCO OD TO	

\*\*\*\* GAMMA DEVIATE GENERATOR \*\*\*

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C METHODS	NORMAL ARRAY INDEX INCREMENT NORMAL ARRAY INDEX LIMIT NORMAL ARRAY INDEX	ARRAY INDEX LIMIT ARRAY INDEX LIMIT		CALL TO SQRT IN "GO" SET UP	2ND CALL TO SORT IN "GO" SET UP	CALL TO ALOG IN "GO" SETUP	CALLS TO REPLENISH RNARRAY		CALL TO ALOG IN NORMAL SECTION OF "GO"	CALL TO ALDG IN EXPON SECTION OF "GO"	CALL TO EXPONENTIAL GENERATOR IN "GO"		CALL TG ALOG IN METHOD "GF"	FUNCTION CALLS IN METHOD "GS"	
CCNSTANTS FOR AD HOC	DC F14 0	DC F 16.	ARGUMENT LISTS	X°FF	1×<	X - X	DC AL4(IX) DC AL4(RNARRAY)	X	i X		AL4	X - X	XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX	1×<	ORG TEST
***	CHICON3		* * *	ARGLST1	ARGLST2	ARGLST3	ARGL ST4		<b>ARGLST5</b>	<b>ARGLST6</b>	ARGLST7		<b>ARGLST8</b>	ARGLST9	

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